# ON INTERPRETATIONS IN BÜCHI ARITHMETICS 

ALEXANDER ZAPRYAGAEV

Büchi arithmetics $\mathbf{B A}_{n}, n \geq 2$, are the theories $\mathbf{T h}\left(\mathbb{N} ;=,+, V_{n}\right)$ where $V_{n}$ is an additional unary functional symbol such that $V_{n}(x)$ is the largest power of $n$ that divides $x$. By convention, we put $V_{n}(0)=0$. They are a series of natural extensions of Presburger arithmetic $\operatorname{Pr} \mathbf{A}=\operatorname{Th}(\mathbb{N} ;=,+)$ [7] that are also complete and decidable. The theories $\mathbf{B A}_{n}$ were proposed by J. Büchi [2] in order to describe the sets of natural numbers recognizable by finite automata through definability in some arithmetical language.

Let $\operatorname{Digit}_{n}(x, y)$ be the digit corresponding to $n^{y}$ in the $n$-are expansion of $x$. We consider automata over the alphabet $\{0, \ldots, n-1\}^{m}$. Assume that at step $k$, the automaton receives the input $\left(\operatorname{Digit}_{n}\left(x_{1}, k\right), \ldots, \operatorname{Digit}_{n}\left(x_{m}, k\right)\right)$, containing the digits corresponding to $n^{k}$ in the $n$-ary expansion of numbers $\left(x_{1}, \ldots, x_{m}\right)$. We say the automaton accepts the tuple $\left(x_{1}, \ldots, x_{m}\right)$ if it accepts the corresponding sequence of tuples $\left(\operatorname{Digit}_{n}\left(x_{1}, k\right), \ldots, \operatorname{Digit}_{n}\left(x_{m}, k\right)\right)$. Then the following well-known result by V. Bruyère ( 3,4$]$ ) holds:

Proposition 1. Let $\varphi\left(x_{1}, \ldots, x_{m}\right)$ be a $\mathbf{B} \mathbf{A}_{n}$-formula. Then there is an effectively constructed automaton $\mathcal{A}$ such that $\left(a_{1}, \ldots, a_{m}\right)$ is accepted by $\mathcal{A}$ iff $\mathbb{N} \models \varphi\left(a_{1}, \ldots, a_{m}\right)$.

Contrariwise, let $\mathcal{A}$ be a finite automaton working on $m$-tuples of $n$-ary natural numbers. Then there is an effectively constructed $\mathbf{B} \mathbf{A}_{n}$-formula $\varphi\left(x_{1}, \ldots, x_{m}\right)$ such that $\mathbb{N} \models \varphi\left(a_{1}, \ldots, a_{m}\right)$ iff $\left(a_{1}, \ldots, a_{m}\right)$ is accepted by $\mathcal{A}$.

Let $\mathbf{T}$ and $\mathbf{U}$ be two first-order theories. We call an $m_{1}$-dimensional interpretation $\iota_{1}$ and $m_{2}$ dimensional interpretation $\iota_{2}$ from $\mathbf{T}$ to $\mathbf{U}$ provably isomorphic, if in the language of $\mathbf{U}$ there is a formula $F(\bar{x}, \bar{y})$ expressing an isomorphism $f$ between the corresponding internal models $\mathfrak{A}_{1}$ and $\mathfrak{A}_{2}$, and the statement that $f$ is an isomorphism must be provable in $\mathbf{U}$.

Note that two interpretations into the elementary theory of a model $\mathfrak{B}$ are provably isomorphic iff there is an isomorphism between their internal models in $\mathfrak{B}$ expressible in the language of $\mathfrak{B}$. As $\mathbf{B} \mathbf{A}_{n}$ is the elementary theory of the model $\left(\mathbb{N} ;=,+, V_{n}\right)$ respectively, it is sufficient to study the interpretations in this model when studying interpretations up to a provable isomorphism.
A. Visser has asked the following question: given an weak arithmetical theory $\mathbf{T}$ without ability to encode syntax but with full induction, does it hold that each interpretation (one-dimensional or multidimensional) of $\mathbf{T}$ into itself is provably isomorphic to the trivial one? This question was previously answered positively for Presburger arithmetic $\operatorname{PrA}$ in the author's joint work 6 with F. Pakhomov.

We show that
Lemma 2. Each $\mathbf{B A}_{k}$ is interpretable in any of $\mathbf{B} \mathbf{A}_{l}, k, l \geq 2$.
Hence, it is sufficient to consider a particular Büchi arithmetic, such as $\mathbf{B A}_{2}$.
In the talk, we establish that each interpretation of $\mathbf{B} \mathbf{A}_{n}$ in itself is isomorphic to the trivial one [8. Furthermore, we show this result holds already for the interpretations of Presburger arithmetic in $\mathbf{B A}_{n}$ :

Theorem 3. Let $\iota$ be a (one- or multi-dimensional) interpretation of $\operatorname{Pr} \mathbf{A}$ in $\left(\mathbb{N} ;=,+, V_{n}\right)$. Then the internal model induced by ८ is isomorphic to the standard one.

From Proposition 1 it follows that an algebraic structure is interpretable in $\mathbf{B} \mathbf{A}_{n}$ iff it has an automatic [5] presentation. Hence, in automatic terms, the statement proved implies there is no automatic non-standard model of any Büchi arithmetic.

The proof is based on the contradiction between the Kemeny-style description of the order types of non-standard models of Büchi arithmetics and the following condition on automatic torsion-free abelian groups established by Braun and Strüngmann [1:

Proposition 4. Let $(A,+)$ be an automatic torsion-free abelian group. Then there exists a subgroup $B$ of $A$ isomorphic to $\mathbb{Z}^{m}$ for some natural $m$ such that the orders of the elements in $C=A / B$ are only divisible by a finite number of different primes $p_{1}, \ldots, p_{s}$.

This gives a partial positive answer to the question of Visser.
Whether the isomorphism of Theorem 3 is always definable by a $\mathbf{B} \mathbf{A}_{n}$-formula remains a problem for future research.
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National Research University Higher School of Economics, Moscow, Russia
Email address: azapryagaev@hse.ru

