Programming BSP and Multi-BSP algorithms in ML

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Structured parallel computing BSP and BSML MULTI-BSP and MULTI-ML

2 Comparison



The world of parallel computing





Distributed computing

Characterised by:

- Interconnected units
- Distributed memory
- Communication network
- MPI



Bulk Synchronous Parallelism

The BSP computer

Defined by:

- p pairs CPU/memory
- Communication network
- Synchronisation unit
- Super-steps execution

Properties:

- Deadlock-free
- Predictable performances



local computations

communication

barrier next super-step

Bulk Synchronous ML

What is BSML?

• Explicit BSP programming with a functional approach



Bulk Synchronous ML

What is $\ensuremath{\operatorname{BSML}}\xspace?$

- Explicit BSP programming with a functional approach
- Based upon ML and implemented over OCAML



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- Explicit BSP programming with a functional approach
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- Formal semantics → computer-assisted proofs (COQ)



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What is BSML?

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Main idea

Parallel data structure \Rightarrow *parallel vector*:





Hierarchical architectures

Characterised by:

- Interconnected units
- Both shared and distributed memories
- Hierarchical memories



MULTI-BSP

- 1 A tree structure with nested components
- 2 Where nodes have a storage capacity
- 3 And leaves are processors
- **4** With sub-synchronisation capabilities



MULTI-BSP

- Stage 3: 4 nodes with a network access
- Stage 2: one node has 4 chips plus RAM
- Stage 1: one chip has 8 cores plus L3 cache
- Stage 0: one core with L1/L2 caches



The $\operatorname{MULTI-BSP}$ model

Execution model

A level *i* superstep is:



The MULTI-BSP model

Execution model

A level *i* superstep is:

• Level i-1 executes code independently



The MULTI-BSP model

Execution model

A level *i* superstep is:

- Level i 1 executes code independently
- Exchanges information with the m_i memory



The $\operatorname{MULTI-BSP}$ model

Execution model

A level *i* superstep is:

- Level i-1 executes code independently
- Exchanges information with the m_i memory
- Synchronises



Basic ideas



Basic ideas

- BSML-like code on every stage of the MULTI-BSP architecture
- Specific syntax over ML: eases programming



Basic ideas

- BSML-like code on every stage of the MULTI-BSP architecture
- Specific syntax over ML: eases programming
- Multi-functions that recursively go through the MULTI-BSP tree



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2 Comparison The methodology Application cases

3 Conclusion

How to compare ? I

"Incremental programming"

- **1** Write a sequential OCAML code
- **2** BSP extension using BSML
- **3** MULTI-BSP extension using MULTI-ML

How to compare ? II

Difficulty of writing a code

- Halsted Difficulty (HD)
- Halsted Effort (HE)
- McCabe Cyclomatic Complexity (CC)
- Maintainability Index (MI)

How to compare ? III

Overall performances

- Speedup (based on sequential algorithm)
- Execution time

Execution platforms

- Mirev2:
 - 8 nodes with 2 quad-cores (AMD 2376 at 2.3Ghz)
 - 16GB of memory per node
 - 1Gbit/s network
- Mirev3:
 - 4 nodes with 2 octo-cores (INTEL xeon E5 2650 at 2.6Ghz)
 - 64GB of memory per node
 - 10Gbit/s network

Application cases

- Model-checking (Symbolic computation)
- Skeleton based FFT (Numerical computation)
- All Pairs similarity search problem (*Big-data computation*)

Model-checking case

What is it ?

Exhaustively check if a model meets a given specification.

How ?

Exploring all the states accessible via a successor function from an initial state ${\it s}_0$

Algorithm characteristics

- Data intensive
- Task parallel
- Explicit load balancing

Difficulty of writing the code



Overall performances



Skeleton based FFT case

What is it ?

Converts a signal from its original domain to a representation in the frequency domain

How ?

 $(FFTx)_i = \sum_{k=0}^{n-1} x_k \cdot e^{2\pi \sqrt{-1}/n_n^{ki}}$ Expressed using the Distributable Homomorphism (DH) skeleton

Algorithm characteristics

- Data intensive
- Data parallel
- Balanced communication scheme

Difficulty of writing the code



Overall performances



V. Allombert, F. Gava

All Pairs similarity search problem case

What is it ?

Discover all the pairs of objects whose similarity is above a given threshold

How ?

Compute local similarities of data sub-sets, then exchanges

Algorithm characteristics

- Data intensive
- Data parallel
- Balanced communication scheme

Difficulty of writing the code



Overall performances



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2 Comparison



To sum up

Model-checking

- Code complexity: OCAML < BSML < MULTI-ML
- Overall performances: OCAML < BSML ≤ MULTI-ML

Skeleton based FFT

- Code complexity: $OCAML < BSML \simeq MULTI-ML$
- Overall performances: $OCAML < BSML \simeq MULTI-ML$

All Pairs similarity search problem

- Code complexity: OCAML < BSML < MULTI-ML
- Overall performances: OCAML < BSML < MULTI-ML

Conclusion

Hierarchical programming: is it worth it ?

+ Performances

- Hard to program
- -> Harder to write *immortal* algorithms

Ongoing and future work

- Write MULTI-BSP algorithms
- Comparison with BSP and MULTI-BSP cost model
- Programming experiments

Thank you for your attention O

 ${\sf Questions}\ ?$