MULTI-ML: PROGRAMMING MULTI-BSP ALGORITHMS IN ML

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BSPSP October 2015





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- 3 Results
- 4 Conclusion

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- 1 Introduction
 OCAML
 BSML
 MULTI-BSP
- 2 Multi-ML
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Strengths of Ocaml

A functionnal programming language

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Strengths of Ocaml

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- A powerful type system



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Strengths of Ocaml

- A functionnal programming language
- A powerful type system
- User-definable algebraic data types and pattern matching
- Automatic memory management
- Efficient native code compilers

Syntaxe overview

```
# let f = fun \times \rightarrow "Hello_"^(string_of_int x) in
  let lst = [0;1;2;3] in
  List.map f lst::
-: string list = ["Hello_0"; "Hello_1"; "Hello_2"; "Hello_3"]
# let pair = ([0;1;2;3],true);;
val pair : int list * bool = ([0; 1; 2; 3], true)
# type 'a list =
   Nil
      Node of 'a*'a list ;;
type 'a list = Nil | Node of 'a * 'a list
```

Bulk Synchronous ML

What is BSML?



Bulk Synchronous ML

What is BSML?

• Explicit BSP programming with a functional approach



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- Based upon ML an implemented over OCAML



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- Formal sementics → computer-assisted proofs (COQ)

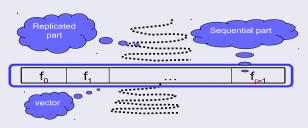


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- Explicit BSP programming with a functional approach
- Based upon ML an implemented over OCAML
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Main idea

Parallel data structure \Rightarrow vectors:



Asynchronous primitives

Asynchronous primitives

•
$$\ll$$
 e \gg : $\langle e, \dots, e \rangle$

Asynchronous primitives

- \ll e \gg : \langle e,..., e \rangle
- \$v\$: v_i on processor i, assumes $v \equiv \langle v_0, \dots, v_{p-1} \rangle$

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Synchronous primitives

• proj : $\langle x_0, \dots, x_{p-1} \rangle \mapsto (\text{fun } i \to x_i)$

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Synchronous primitives

- proj : $\langle x_0, \dots, x_{p-1} \rangle \mapsto (\text{fun } i \to x_i)$
- put : $\langle f_0, \dots, f_{p-1} \rangle \mapsto \langle (\text{fun } i \rightarrow f_i \ 0), \dots, (\text{fun } i \rightarrow f_i \ (p-1)) \rangle$

For a BSP machine with 3 processors:

```
# let vec = \ll "Hello_" >> ;;

val vec : string par = <"Hello_", "Hello_", "Hello_">

# let vec2 = \ll $vec$^(string_of_int $pid$) >> ;;

val vec2 : string par = <"Hello_0", "Hello_1", "Hello_2">

# let totex v = List.map (proj v) procs;;

val totex : 'a Bsml.par \rightarrow 'a list = <fun>

# totex vec2;;

— : string list = ["Hello0"; "Hello1"; "Hello2"]
```

What is MULTI-BSP?

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1 A tree structure with nested components

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- 1 A tree structure with nested components
- 2 Where nodes have a storage capacity

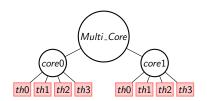
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- 3 And leaves are processors

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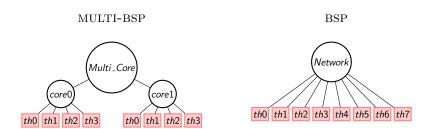
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MULTI-BSP



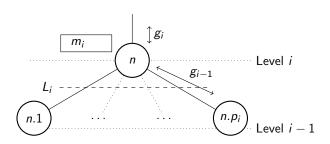
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Execution model

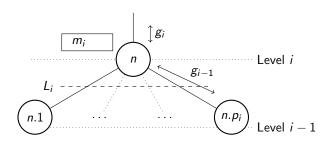
A level *i* superstep is:



Execution model

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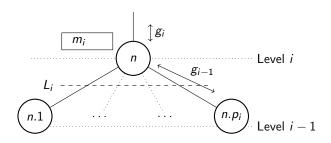
• Level i-1 executes code independently



Execution model

A level i superstep is:

- Level i-1 executes code independently
- Exchanges informations with the m_i memory



Execution model

A level *i* superstep is:

- Level i-1 executes code independently
- Exchanges informations with the m_i memory
- Synchronises

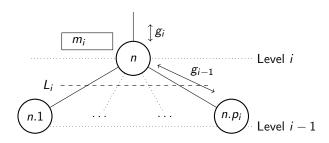


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Overview

Primitives

Semantics

Typing

Implementation

- Results
- 4 Conclusion

Basic ideas:		

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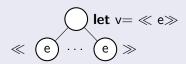
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- BSML-like code on every stage of the MULTI-BSP architecture
- Specific syntax over ML: eases programming
- Multi-functions that recursively go through the tree.

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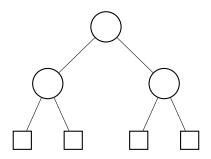


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let multi f [args] =
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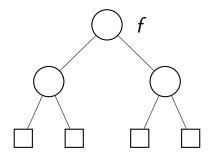
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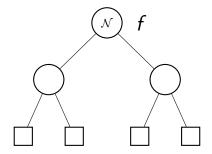


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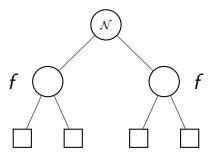
« f [args] »
... in v
where leaf =
(* OCAML code *)
... in v
```



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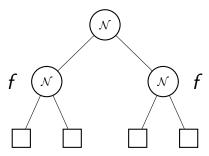


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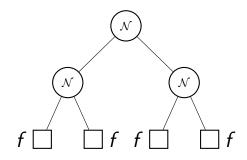
Recursion structure

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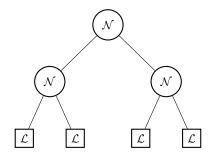
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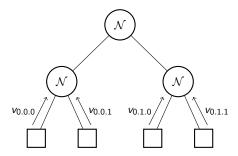


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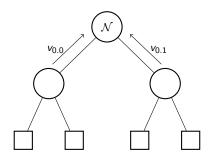


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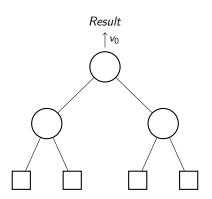


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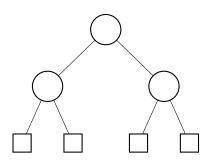


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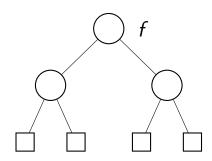


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let multi tree f [args] = where node = (* \operatorname{BSML} \operatorname{code} *) ... in (\ll \operatorname{f [args]} \gg , \operatorname{v}) where leaf = (* \operatorname{OCAML} \operatorname{code} *) ... in \operatorname{v}
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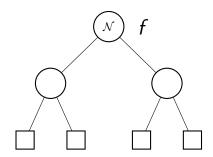
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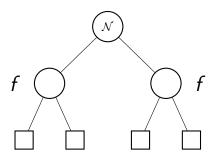
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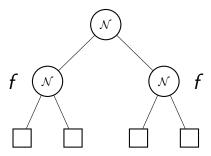
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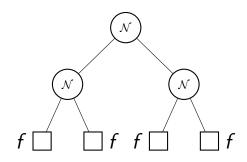
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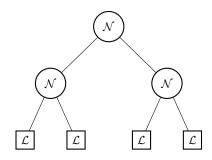
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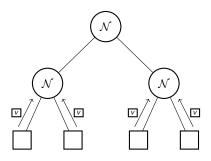
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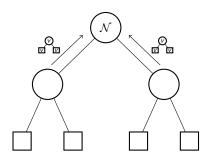
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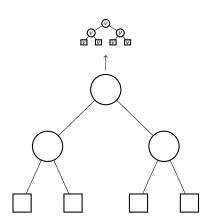
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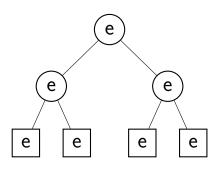


Summary:

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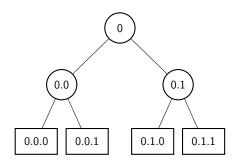
• §e§



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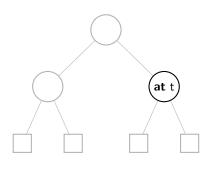
Summary:

- §e§
- gid

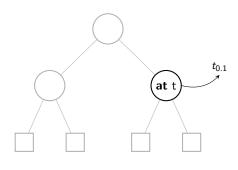


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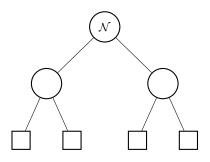
- §e§
- gid
- at



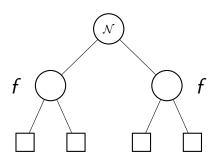
- §e§
- gid
- at



- §e§
- gid
- at
- «...f...»

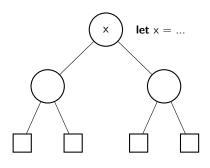


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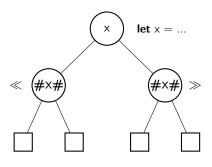
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- §e§
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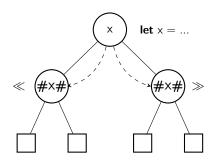


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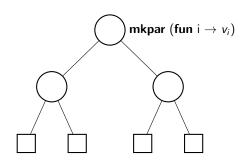
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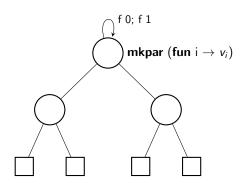
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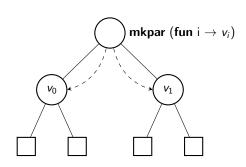
- §e§
- gid
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- «...f...»
- #x#
- mkpar f



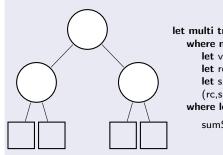
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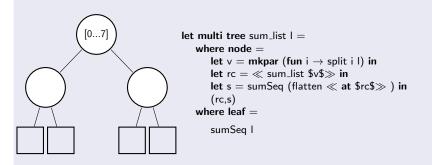


Keep the intermediate results of the sum:

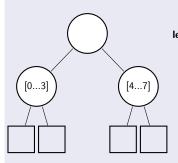


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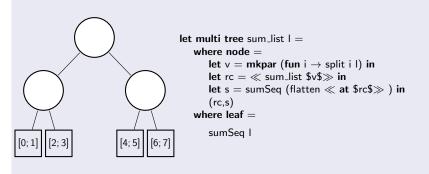


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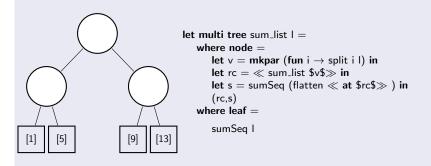


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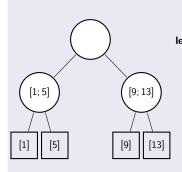
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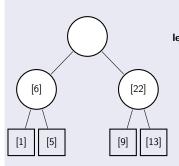


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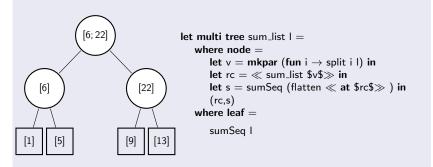
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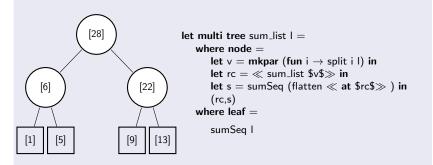


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Semantics

Formal definition of a core-language

Useful for:

- Study of properties
- Proof of programs/compiler/typing rules

Currently

• Inductive big-step: confluent

• Co-inductive: mutually exclusive

Typing

Purely Constraint-Based system : PCB(X)

- Constraint based
- Extension of DM's type system
- Easy to extend
- Related to HM(X)

MULTI-ML type extension

- Add parallel constructions
- Introduce locality $(s, \ell, b \text{ and } m)$ using effects
- Reject nested vectors
- Concistency

Implementation

Sequential simulator

- OCAML-like toplevel
- Test and debug
- Tree sturcture
- Hash tables to represent memories

```
#let multi tree f n =
   where node =
      let r = \ll f ( pid + \#n\# + 1) \gg in
         (r,(gid^"=>"^n))
   where leaf=
      (gid^"=>"^n);;
— : val f : int\rightarrowstring tree = <multi-fun>
\#(f 0)
0 " 0 \rightarrow 0"
--o "0.0→ 1"
\begin{array}{c|c} -\rightarrow & "0.0.0 \rightarrow 2" \\ -\rightarrow & "0.0.1 \rightarrow 3" \end{array}
```

Our approach

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Modular

Our approach

- Modular
- Generic functors

Our approach

- Modular
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- Communication routines

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- Modular
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- Portable on shared and distributed memories

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Current version

Based on MPI

Our approach

- Modular
- Generic functors
- Communication routines
- Portable on shared and distributed memories

Current version

- Based on MPI
- SPMD

Our approach

- Modular
- Generic functors
- Communication routines
- Portable on shared and distributed memories

Current version

- Based on MPI
- SPMD
- One process for each nodes/leaves

Our approach

- Modular
- Generic functors
- Communication routines
- Portable on shared and distributed memories

Current version

- Based on MPI
- SPMD
- One process for each nodes/leaves
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Benchmarks

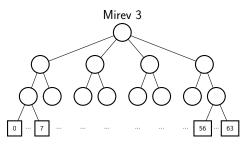
Naive Eratosthenes algorithm

- $\sqrt{(n)}$ th first prime numbers
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Benchmarks

Naive Eratosthenes algorithm

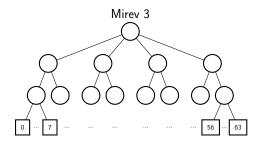
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Benchmarks

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Results

	100_000		500_000		1_000_000		
	MULTI-ML	BSML	MULTI-ML	BSML	MULTI-ML	BSML	
8	0.7	1.8	22.4	105.0	125.3	430.7	
64	0.3	0.3	1.3	8.7	4.1	56.1	
128	0.5	0.45	2.1	5.2	4.7	24.3	

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MULTI-ML

• Recursive multi-functions

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Current/Future work

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Current/Future work

- Optimise MPI implementation
- Type system for MULTI-ML
- Real life benchmarks

Merci!

Any questions ?