

# Higher randomness and forcing with closed sets

Benoit Monin - LIAFA - University of Paris VII



STACS - 07 March 2014

# Overview

The goal of this work is to get a better understanding of two notions of randomness :

- 1/ The first one is called weakly- $\Pi_1^1$ -randomness (Notion introduced by Nies, in 'Computability and randomness', 2009)
- 2/ The second one is called  $\Pi_1^1$ -randomness (Notion introduced by Sacks, in 'Higher recursion theory', 1990)

The goal : Show that 2/ is stronger than 1/

How : Yu and Chong proposed to show that the two notions have different Borel complexity.

We showed that this is not possible, as both sets are  $\Sigma_3^0$ .

# Complexity of sets



## Section 1

# Complexity of sets

# The Cantor space

What do we work with ?

<b>Our playground</b>	The Cantor space
<b>Denoted by</b>	$2^\omega$
<b>Topology</b>	The one generated by the cylinders $[\sigma]$ , the set of sequences extended $\sigma$ , for every string $\sigma$
<b>An open set <math>\mathcal{U}</math> is</b>	A union of cylinders

# Arithmetical complexity of sets

## The Borel sets on the Cantor space

$\Sigma_1^0$ sets are	Open sets
$\Pi_1^0$ sets are	Closed sets
$\Sigma_{n+1}^0$ sets are	Countable unions of $\Pi_n^0$ sets
$\Pi_{n+1}^0$ sets are	Complements of $\Sigma_{n+1}^0$ sets

What about a union of  $\Pi_n^0$  sets for  $n$  unbounded in  $\omega$ ?

# Arithmetical complexity of sets

We define Borel sets by induction over the ordinals

$\Sigma_1^0$ sets are	Open sets
$\Pi_1^0$ sets are	Closed sets
$\Sigma_{\alpha+1}^0$ sets are	Countable unions of $\Pi_{\alpha}^0$ sets
$\Pi_{\alpha+1}^0$ sets are	Complements of $\Sigma_{\alpha+1}^0$ sets
$\Sigma_{\sup_n \alpha_n}^0$ sets are	Countable unions of $\Pi_{\beta}^0$ sets for $\beta < \sup_n \alpha_n$

# Effectivize the arithmetical complexity of sets

$$\text{Notation : } [W_e] = \bigcup_{\sigma \in W_e} [\sigma]$$

$\Sigma_1^0$ sets are	of the form $[W_e]$
$\Pi_1^0$ sets are	of the form $[W_e]^c$
$\Sigma_2^0$ sets are	of the form $\bigcup_{n \in W_e} [W_n]^c$
$\Pi_2^0$ sets are	of the form $\bigcap_{n \in W_e} [W_n]$
...	...

Effectivize  $\Sigma_\alpha^0$  for  $\alpha$  not 'too big' is also doable.

# Randomness and genericity



## Section 2

# Randomness and genericity



# Algorithmic randomness

The general idea of randomness :

- 1 Fix a countable class of sets of measure 1
- 2 Say that something is random if it belongs to all of them

1-random	every $\Sigma_2^0$ sets 'effectively of measure 1'
weakly-2-random	every $\Sigma_2^0$ sets of measure 1
2-random	every $\Sigma_3^0$ sets 'effectively of measure 1'
...	...

# Algorithmic randomness

## Effectively of measure 1 sets

A union of set  $\bigcup \mathcal{A}_n$  is said to be effectively of measure 1 if there is a computable function  $f : \omega \mapsto [0, 1]$  going to 1, such that  $\lambda(\mathcal{A}_n) > f(n)$ .

We have :

1-random  $\leftarrow$  w2-random  $\leftarrow$  2-random  $\leftarrow$  w3-random  $\leftarrow$  ...

All implications are strict

# Genericity

The general idea of genericity :

- ① Fix a countable class of dense open sets of  $2^\omega$
- ② Say that something is generic if it belongs to all of them

weakly-1-generic	every dense $\Sigma_1^0$ set
1-generic	every set $\mathcal{U} \cup (\overbrace{2^\omega - \mathcal{U}}^c)$ for $\mathcal{U}$ a $\Sigma_1^0$ set
weakly-2-generic	every dense $\Sigma_1^0(\emptyset')$ set

weakly-1-generic  $\leftarrow$  1-generic  $\leftarrow$  weakly-2-generic  $\leftarrow$  2-generic  $\leftarrow$   
 $\dots$

All implications are strict

## Another type of genericity : $\Pi_1^0$ -Genericity

We consider genericity for the **topology generated by  $\Pi_1^0$  sets**

### Dense open sets

An open set is a countable union of  $\Pi_1^0$  set. It is **dense** if it intersects any non empty  $\Pi_1^0$  sets.

weakly- $\Pi_1^0$ -generic	every dense $\Sigma_2^0$ set
$\Pi_1^0$ -generic	every set of the form $\mathcal{F} \cup (\overbrace{2^\omega - \mathcal{U}}^{\circ})$ for $\mathcal{F}$ a $\Sigma_2^0$ set

# Mix genericity and randomness : $\Pi_1^0$ -Solovay-Genericity

We still consider genericity for the **topology generated by  $\Pi_1^0$  sets**

## Open sets

An open set is now **'dense'** if it intersects with positive measure any  $\Pi_1^0$  sets of positive measure.

weakly- $\Pi_1^0$ -Solovay-generic	every 'dense' $\Sigma_2^0$ set
$\Pi_1^0$ -Solovay-generic	every $\Sigma_2^0$ set $\mathcal{F}$ together with all $\Pi_1^0$ sets of positive measure, included in $2^\omega - \mathcal{F}$ .

# Results on $\Pi_1^0$ -Solovay-Genericity

## Definition

A sequence  $X$  is **computably dominated** if any functions  $f : \omega \rightarrow \omega$  that  $X$  can compute is dominated by a computable function  $g : \omega \rightarrow \omega$  (that is we have for all  $n$  that  $f(n) \leq g(n)$ ).

## Theorem

A sequence  $X$  is  $\Pi_1^0$ -generic iff it is computably dominated.

## Theorem

A sequence  $X$  is  $\Pi_1^0$ -Solovay generic iff it is computably dominated and weakly-2-random.

# Beyond arithmetic



## Section 3

Beyond arithmetic

# Computable ordinals

## Computable ordinals

An ordinal  $\alpha$  is computable if there is a c.e. well-order  $R \subseteq \omega \times \omega$  so that  $|R| = \alpha$ .

## $\mathbf{W}$ and $\mathbf{W}_\alpha$

We denote by  $\mathbf{W}$  the set of codes for computable ordinals, and we denote by  $\mathbf{W}_\alpha$  the set of codes for computable ordinals strictly smaller than  $\alpha$ .

## Strict initial segment

The computable ordinals forms a strict initial segment of the countable ordinals.



# Analytic and co-analytical sets

## $\Pi_1^1$ sets

A set of integers/sequence is  $\Pi_1^1$  if it definable by a formula of arithmetic whose quantifiers over sets are only universal.

## $\Sigma_1^1$ sets

A set of integers/sequence is  $\Sigma_1^1$  if it definable by a formula of arithmetic whose quantifiers over sets are only existential.

## $\Delta_1^1$ sets

A set of integers/sequence is  $\Delta_1^1$  if it both  $\Pi_1^1$  and  $\Sigma_1^1$ .

# Analytic and co-analytical sets

## Theorem

A set of integers  $A$  is  $\Pi_1^1$  iff there is a computable function  $f : \omega \rightarrow \omega$  so that  $n \in A$  iff  $f(n) \in \mathbf{W}$ .

A set of sequences  $\mathcal{A}$  is  $\Pi_1^1$  iff there is a code  $e$  so that  $X \in \mathcal{A}$  iff  $e \in \mathbf{W}^X$ .

## Theorem

A set of integers  $A$  is  $\Delta_1^1$  iff there is a computable function  $f : \omega \rightarrow \omega$  and a computable ordinal  $\alpha$  so that  $n \in A$  iff  $f(n) \in \mathbf{W}_\alpha$ .

A set of sequences  $\mathcal{A}$  is  $\Delta_1^1$  iff there is a code  $e$  and a computable ordinal  $\alpha$  so that  $X \in \mathcal{A}$  iff  $e \in \mathbf{W}_\alpha^X$ .

## Example of $\Pi_1^1$ set of sequence

 $\omega_1^{ck}$ 

We denote by  $\omega_1^{ck}$  the smallest non-computable ordinal.

 $\omega_1^X$ 

We denote by  $\omega_1^X$  the smallest non-computable-in- $X$  ordinal.

The set  $\{X : \omega_1^X > \omega_1^{ck}\}$  is :

- A  $\Pi_1^1$  set
- of measure 0
- meager

# Higher randomness/genericity



## Section 4

Higher randomness &  
genericity

## Another hierarchy

### $\Pi_1^1$ open sets & $\Sigma_1^1$ closed sets

An open set is a  $\Pi_1^1$  open set if it can be described with a  $\Pi_1^1$  set of strings. The  $\Sigma_1^1$  closed sets are complements of  $\Pi_1^1$  open set.

We can have a 'higher Borel hierarchy' starting from  $\Pi_1^1$  open sets and  $\Sigma_1^1$  closed sets :

- 1 A set is  $\Pi_2^0(\Pi_1^1)$  if it is a uniform intersection of  $\Pi_1^1$  open sets.
- 2 A set is  $\Sigma_2^0(\Sigma_1^1)$  if it is a uniform union of  $\Sigma_1^1$  closed sets.
- 3 A set is  $\Pi_3^0(\Sigma_1^1)$  if it is a uniform intersection of a uniform union of  $\Sigma_1^1$  closed sets.
- 4 A set is  $\Sigma_3^0(\Pi_1^1)$  if it is a uniform union of a uniform intersection of  $\Pi_1^1$  open sets.

# Higher randomness

## $\Delta_1^1$ -random

A sequence is weakly- $\Delta_1^1$ -random if it belongs to no  $\Delta_1^1$  nullset

## weakly- $\Pi_1^1$ -random

A sequence is weakly- $\Pi_1^1$ -random if it belongs to no  $\Pi_2^0(\Pi_1^1)$  nullset

## $\Pi_1^1$ -random

A sequence is  $\Pi_1^1$ -random if it belongs to no  $\Pi_1^1$  nullset

We have

$$\Delta_1^1\text{-random} \leftarrow \text{weakly-}\Pi_1^1\text{-random} \leftarrow \Pi_1^1\text{-random.}$$

# Higher randomness

Is  $\Pi_1^1$ -random different from weakly- $\Pi_1^1$ -random ?

An approach to solve this question is to show that the set of  $\Pi_1^1$ -randoms and the set of weakly- $\Pi_1^1$ -randoms have different Borel complexity. We showed that is is not possible, as both sets are  $\Sigma_3^0$ .

## Higher randomness with genericity

We give an analogue of the notions of Solovay genericity with  $\Pi_1^0$  sets.

### weakly- $\Sigma_1^1$ -Solovay-generic

A sequence is weakly- $\Sigma_1^1$ -Solovay-generic if it is in every  $\Sigma_2^0(\Sigma_1^1)$  sets intersecting with positive measure every  $\Sigma_1^1$  closed set.

### $\Sigma_1^1$ -Solovay-generic

A sequence  $X$  is  $\Sigma_1^1$ -Solovay-generic if for every  $\Sigma_2^0(\Sigma_1^1)$  sets  $\mathcal{F}$ , either  $X$  is in it or  $X$  belongs to a  $\Sigma_1^1$  set of positive measure, included in the complement of  $\mathcal{F}$ .



# Equivalence 1

## Fact

The set of weakly- $\Pi_1^1$ -randoms is obviously  $\Sigma_3^0$ .

## Theorem 1

A sequence is weakly- $\Sigma_1^1$ -Solovay-generic iff it is weakly- $\Pi_1^1$ -random.

- If  $X$  is weakly- $\Sigma_1^1$ -Solovay-generic then it is surely in every  $\Sigma_2^0(\Sigma_1^1)$  sets of measure 1.
- ← If  $X$  is not weakly- $\Sigma_1^1$ -Solovay-generic then it is in a  $\Pi_2^0(\Pi_1^1)$  containing no  $\Sigma_1^1$  closed set of positive measure. But such a  $\Pi_2^0(\Pi_1^1)$  is necessarily of measure 0 (Not trivial).

## Equivalence 2

### Theorem 2

A sequence is  $\Sigma_1^1$ -Solovay-generic iff it is  $\Pi_1^1$ -random

### Theorem Chong, Yu, Nies

A sequence is  $\Pi_1^1$ -random iff it is  $\Delta_1^1$ -random and  $\omega_1^X = \omega_1^{ck}$ .

Proof Idea of theorem 2 :

- ← If  $X$  is not  $\Sigma_1^1$ -Solovay-generic then it is in a  $\Pi_2^0(\Pi_1^1)$  set  $\mathcal{G}$  so that for no  $\Sigma_1^1$  closed set of positive measure  $\mathcal{F}$  we have  $X \in \mathcal{F} \subseteq \mathcal{G}$ . If  $\omega_1^X = \omega_1^{ck}$  we can deduce from this that  $X$  is in a  $\Delta_1^1$  set of measure 0 (not trivial).
- If  $X$  is  $\Sigma_1^1$ -Solovay-generic then surely  $X$  is  $\Delta_1^1$ -random. All that need to be proved is that  $\omega_1^X = \omega_1^{ck}$ . This is the tricky part.

# Consequences

## Corollary

The set of  $\Pi_1^1$ -randoms has the same Borel complexity than the weakly- $\Pi_1^1$ -randoms :  $\Sigma_3^0$ . Using a result of Liang Yu, this complexity is strict (it is not  $\Pi_3^0$ ).

However, using a another technique, we recently proved that :

## Last results

There is some  $X$  which is weakly- $\Pi_1^1$ -random but not  $\Pi_1^1$ -random.

Thank you