

Normalization of AsmL programs

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15-06-09



OUTLINE

- Introduction: *ASM versus AsmL*
- Definitions: syntax and semantics
- Normalization



ASM

ASM

- Algorithmically complete language
- Purely functional
- Purely parallel



AsmL

AsmL

- Implement ASM
- Closer to usual languages
- Easier to program with



ASM versus AsmL

ASM:

basic constructs (rules)

- Updates
- Tests
- Parallel composition

More Data Structures

AsmL:

basic constructs (rules)

ASM plus

- sequential composition
- UNTIL and WHILE loops

More Control Structures

ASM Syntax

ASM Syntax: rules are

- Updates: $f(t_1, \dots, t_n) := t_0$
- Tests: IF φ THEN R_1 ELSE R_2 ENDIF
- Parallel composition:

```
par
  R1
  ⋮
  Rk
endpar
```

ASM Semantics

ASM Semantics: rules are executed in parallel until a fixed point is reached

Example: compute $\text{gcd}(a,b)$

IF ($a \bmod b = 0$) THEN $g := b$ ELSE

par

$a := b$

$b := a \bmod b$

endpar

ENDIF

AsmL Syntax

AsmL Syntax: rules are

- ASM rules
- Sequential composition: **step rule**

```
par  
  step  $R_1$   
   $\vdots$   
  step  $R_k$   
endpar
```

- **Iterations:**

```
step until  $\varphi$   $R$   
step while  $\varphi$   $R$   
step until fixpoint  $R$ 
```



AsmL Semantics

AsmL Semantics:

As usual, The **step rule** is interpreted as the **sequential** composition of R_1 , followed by R_2 , etc. , ending with R_k .

par

step R_1

\vdots

step R_k

endpar



AsmL Semantics: example

step

avg := 0

i := 0

step until (i = n)

avg := avg + grade[i]

i := i+1

step

avg := avg/n

i := 0

nb := 0

step until (i = n)

if (grade[i] > avg) then nb := nb+1

i := i+1

Compute the number of grades whose value is greater than the average grade.



Normalization

Goal

Translate the control structures of AsmL into ASM

- Simulate sequential composition.
- Simulate nested iterations with single fixed point iteration.



Normalization: example

```
if (mode = 0) then  avg := 0
                   i := 0
                   mode := 1
if (mode = 1) then  avg := avg + grade[i]
                   i := i+1
                   if (i = n) then mode := 2
if (mode = 2) then  avg := avg/n
                   i := 0
                   nb := 0
                   mode := 3
if (mode = 3) then  if (grade[i] > avg) then nb := nb+1
                   i := i+1
                   if (i = n) then mode := 4
```



Normalization

In the normal form

- *sequential composition* is simulated by controlling the parallelism via a new constant symbol *mode*.
- iteration *until fixpoint* will be simulated by a new control constant *c* which is set to 1 if non-trivial updates are performed (0 otherwise).



A Normal form AsmL program is a program of the form

```
step  until  fixpoint
      par
          if (mode = 0) then  $R_1$ 
          ...
          if (mode = j) then  $R_j$ 
          ...
          if (c=0)  $\wedge$  (mode = i) then  $R_i$ 
          if (c=1)  $\wedge$  ( $M_i \geq$  mode  $\geq$  i) then
              mode:=i
              c:=0
          ...
          if (mode = k) then  $R_k$ 
      endpar
```



Construction of normal forms

By structural induction

AsmL rule	Normal form
ASM rule R R	if (mode = 0) then R mode := 1
AsmL step rule step R_1 step R_2	par R_1'' $R_2^{+fin_1}$ endpar



Construction of normal forms

Iterations

AsmL rule	Normal form
UNTIL STEP UNTIL φ R	if (mode = 0) then if $\neg\varphi$ then mode := 1 if φ then mode := M+1 R^{+1}
UNTIL FIXPOINT step until fixpoint R	if (c=1) \wedge ($M \geq$ mode \geq 0) then mode:=0 c:=0 if (c=0) \wedge (mode = 0) then R^c



Theorem

For every AsmL program P , there exists a normal form ASM program P_n such that for every \mathcal{L} -state \mathcal{A} there exists an \mathcal{L}_m -state \mathcal{A}_m in which P_n computes the same function as P .

Complexity counted by the number of updates: the complexity of P is of the form $\prod n_i$ and the complexity of P_n is of the form $\prod(1 + n_i) > C \times \prod n_i$



Summary

- for each AsmL program P , there is an ASM program P_n in normal form computing the same function
- if the nesting of iterations is bounded by constant k in P , then the complexity of P_n is at most $2k$ times the complexity of P
- can we do better?

THANKS FOR YOUR ATTENTION

